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# TREE-STRUCTURED INFORMATION FILE AND ITS SUBPROGRAM SUBTREE

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The goal of automatic documentation of computer programs is to establish procedures, called documentation programs, that can be implemented by computer programs. These documentation programs may be divided into two categories: postmortem and developmental documentation programs. In the former case, a computer program is presented as input for documentation without any preparation; in the latter case, the program to be documented must be developed so that it contains information necessary for the documentation.

This paper is concerned only with the development documentation programs. A document tree is defined as the syntactic representation of a document when it is divided into subdivisions such as chapters and sections. A developmental tree is defined as a tree of information obtained during the course of the development of a computer program. The task of documenting a computer program is then made equivalent to a transformation of its developmental tree into a document tree. When this transformation is performed by a computer program, the documentation can be achieved automatically.

There is no attempt made in this paper to define the document tree more precisely. Only its tree structure is assumed. Efforts are concentrated on the developmental tree, specifically a subtree of it; the subprogram tree is illustrated in more detail.

#### **GENERAL APPROACH**

In the development of documentation programs, two objectives are paramount. Pieces of information about the program to be documented should be kept in a computer file during the development of the program, and this information should not be duplicated in the file. The importance of the first objective is obvious; the information should be in a computer-readable form for documentation. The importance of the second objective can be seen whenever a change is made during or after the development of the program to be documented. One can easily make the mistake of changing information in one place and forgetting about it in the other place. On the other hand, a change of information at a certain place may require changes in other information.

The goal of this project is to structure the developmental file of information in a tree structure (fig. 1) so that the nodes represent pieces of information. Any change in the

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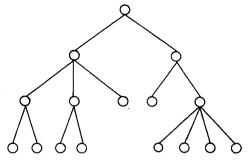


Figure 1.—Tree structure.

contents of a node may require changes in the subtree rooted in that node. In certain cases when the semantic structure is more complex, i.e., it may represent a directed graph, pointers may be used semantically.

The final documentation of a program is produced from its developmental tree of information. A special tree-traversing program, possibly interactive, selects out the contents of nodes or subtrees, invokes certain docu-

mentation programs to transform these data into special format, and stacks this information sequentially. The sequentially stacked information is processed by a listing program to produce the final printed document.

Obviously the main problem is the establishment of the developmental tree structure. At this time, a complete tree structure cannot be proposed. The definition of certain types of subtrees, however, has been accomplished. One of these, a source program subtree, is described in detail.

#### FLOWCHARTING AND PROGRAM LISTINGS

Any large computer program should be segmented into subprograms, subroutines, and procedures. The size of a subprogram may depend on its complexity and on its source language. Documentation of a subprogram is usually done in three different forms: textual description, flowchart, and source language listing.

The information should be structured as a tree. A source program is compiled (assembled), which generates a relocatable program. Figure 2 then defines the tree.

Certain information such as size, entry points, and external references can be obtained from the compiler-generated relocatable program. The rest of the information should be put into the source program. Textual information can easily be placed into the source program by grouped comment lines. Thus the source program may be defined as a tree, as seen in figure 3.

To combine the flowchart with the source program creates some problems. A special



Figure 2.—Tree structure for subprogram.



Figure 3.—Tree structure for source program.

form called a sequence chart is used. This is not a complete flowchart in the standard sense, but it forces a tree on the otherwise graph-structured flowchart. Then there is no problem in listing a tree structure sequentially. The missing links of the graph structure, which appear as transfer statements in the source program, can be implemented by semantic comments. A special computer program for a source language can automatically flag these places.

Appendixes A, B, and C show the final printed forms of three different subprograms. The right side of the lists contains the actual program statements; the left side is stored internally as coded comments. The listing program takes care of this separation, but the actual sequential form is kept in the vertical direction. Those flow lines that represent the spanning tree of the program are shown with special characters, colons, periods, and asterisks. The groups of textual descriptions are separated by horizontal lines of asterisks. Both the names of the groups and the characters used for line drawing are made flexible by changing an internal table in the printing program. Special print programs are available: A "level" print gives only those lines that are not indented more than a certain input parameter. A "selective" print gives only a subtree; i.e., a defined group or a subtree of the body. The output of these print routines, formatted for a document processor, can be kept in the computer.

This form of documentation has been very helpful in the project from which these three examples were taken. During the debugging stage, it was easy to follow the sequence chart to locate a specific segment of a subprogram without turning pages back and forth.

Obviously, to get these forms, a good editing program capable of performing insertions and changes is needed. Appendixes D and E show appendix A in a developmental stage. In appendix D the initial sequence chart is defined. In appendix E an update procedure is shown. First the sequence chart is shown in a coding sheet geometrically; then its code is placed in front of it. The code for a line is composed by two fields. The first field defines either the depth of the text, 0 to 9, and blanks for program statements or contains special instructions, like group heading, change, and insert commands. The second field contains subcodes, such as line drawing codes for sequence charts and line numbers for updating commands. The text appears in the third field. In the actual input, the text field gets left adjusted. The lines will not be represented because they are already defined by codes.

This procedure for writing a program has the following advantages:

- (1) It provides an up-to-date documentation of the program in the developmental stage.
- (2) It forces a programmer to lay out his program so that it provides an automatic documentation at any level.
- (3) It provides a form for a project leader to define subprograms without details that can be inserted by other programmers.
- (4) It may be used for the present-day coded flowcharting programs.

Its main disadvantage is that it needs more work and discipline in the beginning.

# **SUMMARY**

Printed documents have syntactic tree structures, such as titles, chapters, and sections. The semantic contents of the document may have more complex graph structures, but these

structures are implemented by semantic references. A computer program has a graph structure also, but a spanning tree on this graph can be defined with semantic references to the missing links. This developmental tree of a program may have a different arrangement from a document tree. If the necessary information is contained in the developmental tree for the document tree, a transformation program can produce a document tree from the developmental tree. If the structures of the two trees are standardized, then this transformation can be achieved automatically. Otherwise, an interactive transformation routine can achieve a semiautomatic documentation.

#### APPENDIX A-PRINTED SUBPROGRAM: EXAMPLE 1

COUNT=1.

SUBROUTINE EXPRES (\*, ISW) \*\*\*\*\*\*\*\*\*\*\*\*\* TITLE EXPRESSION TRANSLATOR, INFIX TO PREFIX \*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\* ABSTRACT \*\*\* AUTHOR: C.K.MESZTENYI \*\*\* DATE: JULY 21, 1970 ### LAGUAGE: FORTRAN 5 \*\*\* PROJECT: FORMAL - SUBROUTINE \*\*\* SEARCH KEYS: NONE \* DATA STRUCTURE . FORMAL . CMMN INCLUDE CMMN FORMAL . PWORD INCLUDE PWORD \* ERROR RETURN \*\*\* ARGUMENT: ISW INPUT ARGUMENT \*\*\*\*\*\*\*\*\*\*\*\* SPECIFICATION THIS IS A GENERALIZED EXPRESSION TRANSLATION ROUTINE FROM INFIX TO PREFIX FORM. IT ASSUMES THAT THE CALLING ROUTINE INITIALIZED THE SCANNER, THUS GSCANR GIVE THE CONSECUTIVE LOGICAL SYMBOLS. THE ROUTINE MAY BE CALLED FROM 4 DIFFERENT PLACES DEPENDING ON ISW: ISW = 0 PROCESS AN ASSIGN STATEMENT: VARIABLE = EXPRESSION ; = 1 TRANSLATE THE EXPRESSION PART FROM A READ-IN DATA WHICH MAY BE IN THE FORM: EXPRESSION ! OR VARIABLE = EXPRESSION ; = 2 PROCESS SUBSCRIPT EXPRESSION IN THE FORM: EXPRESSION ) = 3 PROCESS AN EXPRESSION IN THE FORM: EXPRESSION / IN THE FIRST CASE, THE INFORMATIONS FOR THE VARIABLE ARE STORED IN NI.NZ.NJ. IN THE SECOND CASE, ONLY THE EXPRESSION PART IS RETAINED UPON RETURN. IN ALL CASES, THE TRANSLATED AND SIMPLIFIED EXPRESSION IS PLACED ABOVE THE PUSH-DOWN STACK WITH THE PUSH-DOWN STACK CONTAINING ONLY ONE ENTRY: A COMMA WITH A COUNT CORRESPONDING THE NUMBER OF EXPRESSIONS TO ACCOMODATE LISTS. \* METHOD AFTER INITIALIZATION, THE LOGICAL BCD SYMBOLS ARE OBTAINED BY GSCANR AND PROCESSED ONE-BY-ONE IN A LOOP. PROCESSING A SYMBOL IS DONE AS FOLLOWS:
FIRST, IT IS CHECKED IF THE SYMBOL IS IN CORRECT TEXT; THEN CONSTANTS- ARE LINKED IN ABOVE THE PUSH-DOWN STACK; VARIABLES - THEIR VALUES ARE OBTAINED FROM THE SYMBOL TABLE AND LINKED ABOVE THE PUSH-DOWN STACK. IF THE VARIABLE IS SUBSCRIPTED, OR IT IS A FUNCTION IDENTIFIER, THEN THE NAME IS LINKED IN ABOVE THE PUSH-DOWN STACK, AND A LEFT PARENTH, IS PLACED IN THE PUSH-DOWN STACK WITH

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LEFT PARENTH, - IS PLACED IN THE PUSH-DOWN
STACK WITH COUNT=0.

OPERATORS - THE PUSH-DOWN STACK IS EMPTIED OUT BY STKOUT
UNTIL ITS TOP ELEMENT HAS PRECEDENCE NUMBER
EQUAL TO OR LESS THAN THE PRECEDENCE NUMBER
OF THE OPERATOR. THEN THE OPERATOR IS PLACED
IN THE PUSH-DOWN STACK. SIMPLIFICATION IS
PERFORMED BY STKOUT.

RIGHT PARENTH., RIGH BRACKET - THE PUSH-DOWN STACK IS
EMPTIED OUT BY STKOUT UNTIL THE MATCHING LEFT
PARENTH. IS FOUND. IF THAT HAS A COUNT=0,
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EMPTIED OUT BY STKOUT UNTIL THE MATCHING LEFT PARENTH. IS FOUND. IF THAT HAS A COUNT=0, IT IS DISCARDED TOGETHER WITH THE RIGHT PARENTH. IF IT HAS A NON-ZERO COUNT, THEN IT INDICATES AN END OF SUBSCRIPTS (PAR.) OR END OF FUNCTION ARGUMENTS (BRACKET). IN CASE OF END OF SUBSCRIPTS, THE SUBSCRIPTS ARE COLLECTED AND THE VALUE OF THE SUBSCRIPTED VARIABLE IS OBTAINED FROM THE SYMBOL TABLE, WHICH IS LINKED IN. IN CASE OF END OF ARGUMENT LIST, THE FUNCTION IDENTIFIER IS OBTAINED AND LINKED IN

SEMICOLON - INDICATES THE END OF EXPRESSION. THE PUSH-DOWN STACK IS EMPTIED OUT BY STKOUT.

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LOGICAL VARIABLE 'SB' IS TRUE WHENEVER THE SCANNED SYMBOL IS IN SUBSCRIPT LEVEL. 'SBC' VARIABLE CONTAINS THE DEPTH OF THIS LEVEL.

LOGICAL VARIABLE 'EQL' IS TRUE WHEN AN '=' HAD BEEN PROCESSED ALREADY, THUS IT MAY NOT APPEAR AGAIN. '=' MAY ALSO NOT APPEAR ON SUBSCRIPT LEVEL.

THE SYNTAX OF EXPRESSIONS IS CHECKED AT EVERY SCANNED SYMBOL BY MASKING 'TEST' WHICH WAS SET BY THE PREVIOUS SYMBOL. IF THE RESULT IS NOT ZERO THEN THE EXPRESSION HAS SYNTACTIC ERROR. IN THE FOLLOWING TABLE, 'A' DENOTES AN ALPHANUMERIC NAME, 'N' DENOTES A NUMERIC CONSTANT,'I' DENOTES POSITIVE INTEGER:

1101151121				
SYMBOL	MASKING BITS	(DEC.)	RESET TEST	(DEC.
INITIAL ASSIG	N		1000000	(64)
INITIAL OTHER	5		0100000	(32)
A	0001110	(14)	0001000	(8)
A (	0001110	(14)	0100000	(32)
AΣ	1001110	(78)	0100000	(32)
N	1001110	(78)	0000100	(4)
[1]	1001110	(78)	0000100	( 4)
# I	1001110	(78)	0000100	(4)
(	1001110	(78)	0100000	(32)
=	1110101 (	117)	0000001	(1)
UNARY +-	1011110	(94)	0010000	(16)
BINARY +-	1110001 (	113)	0010000	(16)
* / **	1110001 (	113)	0010000	(16)
,	1110001 (	(113)	0100000	(32)
) AS SEPARAT	OR 1110001 (	113)	0000100	(4)
3	1110001	113)	0000100	(4)
) AS END OF	SUBS.1110001 (	113)	0000010	( 2)

#### 1110001 (113)

'BRT' AND 'PAR' ARE USED TO COUNT THE BRACKETS AND PARENTHESIS, RESPECTIVELY. LOGICAL 'NEG' IS SET TO TRUE BY '-' FOR THE NEXT CHARACTER SCANNED ONLY.

\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*

BRANCH BY TERMINATING CHARACTER

LOGICAL SB, EQL, NEG

```
SEQUENCE CHART
INITIALIZE
     PUSH-DOWN STACK WITH COMMA
                                                                 NP=IGETF1($990)
                                                                 NP0=NP
                                                                 C(NP)=20K10
                                                                 D(NP)=1
     SUBSCRIPT LEVEL
                                                                 SB= .FALSE.
                                                                 BRT = 0
                                                                 PAR = 0
                                                                 SBC=0
     LOGICAL VARIABLES EQL AND END, INITIAL TEST
                                                                 EQL= ISW .GE. 2
                                                                 TEST=32
                                                                 IF (ISW .EQ. 0) TEST=64
                                                                 NEG= .FALSE.
     GO TO SUBSCRIPT START IF ISW=2
                                                                 IF (ISW .EQ. 2) GO TO 180
LOOP TO PROCESS CONSECUTIVE SYMBOL
     GET SYMBOL
                                                           30
                                                                 CONTINUE
                                                                 CALL GSCANR($990, IND, N1, ITC, ICC)
     BRANCH BY TYPE OF SYMBOL
          IND = 1,2,3,4 FOR
          INTEGER, REAL, IDENTIFIER, SPECIAL CHARACTER
                                                                 GO TO (100,110,40,60), IND
     .... INTEGER
                                                           100
                                                                 I=0
                                                                 GO TO 120
          REAL
                                                           110
                                                                 I=3
          LINK IN CONSTANT
                                                                 IF (AND(TEST, 78) .NE. 0) CALL FMLERR($990, N1, I, 1)
                                                                 TEST=4
                                                                 J=ILINK1(NP,I,N1)
          CHANGE SIGN IF NEG IS TRUE
                                                                 IF (NEG) D(J)=-D(J)
                                                                 NEG= .FALSE.
                                                                 GO TO 30
          IDENTIFIER
          CHECK IS -1 FACTOR SHOULD BE LINKED IN
                                                                 INEG = 1
                                                                 GO TO 500
          ERROR IF IT HAS MORE THAN 6 CHARACTERS
                                                                 IF (ICC .NE. 0) CALL FMLERR($990.N1.1.2)
```

```
GO TO (130: 180: 190): ITC + 1
:... IDENTIFIER NOT TERMINATED BY ( OR [
                                                     IF (AND(TEST,14) .NE. 0) CALL FMLERR($990,N1,1,1)
                                                      TEST=8
                                                      N2=0
     CHECK IF ITS VALUE MUST BE LINKED IN
                                                      IF (EQL .OR. SB) GO TO 160
     :... NO, GET ITS NAME AS VALUE
                                                      IF (N2 .NE. 0) CALL ILINK1(NP.N2+7.N3)
                                                      J=6
                                                      IF (N2 .NE. 0) J=7
                                                      CALL ILINK1 (NP, J, N1)
                                                      GO TO 30
         YES, GET VALUE FROM SYMBOL TABLE
          IF UNASSIGNED, THEN GET
         ITS NAME AS ITS VALUE
                                                      CALL SYMBOL ($990,1)
                                                      IF (EPTR .EQ. 0) GO TO 150
         COPY EXPRESSION AND
         LINK IT WITHOUT LEADING COMMA
                                                      II=ICOPYO($990,EPTR)
                                                      I=NEXT(II)
                                                      J=LASTXX($990,II,1,0)
          IS IT A LIST
                                                      IF (H2(II) .EQ. 1) GO TO 170
         :... YES, EMPTY PUSH-DOWN STACK
               COMBINE COUNT FOR COMMA
                                                      CALL STKOUT ($990,18)
                                                      IF (ITYP(NP) .GT. 17) CALL FMLERR($990,N1,1,1)
                   :
                                                      D(NP)=D(NP)+H2(II)-1
         LINK IN EXPRESSION
                                                      CALL RMOVF1(II)
                                                      CALL ILINKN(NP,I,J)
                                                      GO TO 30
    IDENTIFIER TERMINATED BY LEFT
    PARENTHESIS : A(
                                                      IF (AND(TEST,14) .NE. 0) CALL FMLERR($990,N1,1,1)
                                                      PAR = PAR+1
                                                      TEST=32
     SUBSCRIPTED VARIABLE, LINK IN NAME
     AND PLACE '(' WITH COUNT 1 INTO THE
     STACK. INCREASE SUBSCRIPT LEVEL
                                                      NP=ILINK1(NP,17,1)
                                                      CALL ILINK1 (NP, 7, N1)
                                                      SB= .TRUE.
                                                      SBC=SBC+1
                                                      GO TO 30
:--- IDENTIFIER TERMENATING WITH LEFT
     BRACKET : AE
                                                190 IF (AND(TEST, 78) .NE. 0) CALL FMLERR($990, N1, 1, 1)
     GET FUNCTION IDENTIFIER,
    BRANCH BY TYPE
                                                      I=IFUNCT(N1)
                                                      IF (I .EQ. 0) GO TO 210
                                                      BRT = BRT+1
                                                      IF (1 .GT. 1) GO TO 200
    :... DIFFERENTIAL FUNCTION
```

```
NP=ILINK1(NP,23,0)
                                                       GO TO 240
          MATH. OR FORMAL FUNCTION
                                                      NP=ILINK1(NP+21+0)
                                                 200
                                                       H1 (NP)=I
                                                       GO TO 240
     ... NONE OF ABOVE, CHECK SYMBOL TABLE
                                                 210
                                                       CALL SYMBOL ($990.1)
                                                 220
                                                       BRT = BRT+1
                                                       IF (EPTR .EQ. 0) 60 TO 230
               DEFINED FUNCTION
               LINK IN EXPRESSION
                                                       11=1COPY0($990.EPTR)
                                                       I=NEXT(II)
                                                       J=LASTXX($990,II,1,0)
                                                       NP=ILINK1(NP,22,0(11))
                                                       CALL ILINKN(NP,I,J)
                                                       CALL IFREE1(II)
                                                       GO TO 240
               UNDEFINED FUNCTION
                                                       NP=ILINK1(NP,24,N1)
                                                 230
               LINK IN COMMA FOR
               THE ARGUMENTS FOLLOWING
                                                       NP=ILINK1(NP, 16, 1)
                                                 240
                                                       TEST=32
                                                       GO TO 30
SPECIAL CHARACTERS
BRANCH BY THE SPECIAL CHARACTERS
                                                       GO TO (270,280,290,270,300,340,350,270,440,270,270,270,390,400,
                                                      1 420,270,270,270,270,430,270,460,460,410,270,270,270,N1
:... ILLEGAL SPECIAL CHARACTERS
                                                 270
                                                       CALL FMLERR($990,N1,1,1)
:... LEFT BRACKET OR ID. ENCLOSED IN BRACKETS
      [ [ [ ]
                                                       IF (AND (TEST, 78) .NE. 0) CALL FMLERR ($990, N1, 1, 1)
                                                 280
     UNDEFINED FUNCTION ARGUMENT OR SUBSCRIPTED
     FUNCTION CHECK IF IT IS SUBSCRIPTED
     FUNCTION
     :... YES. GO TO FUNCTION PART
          OF DEFINITION
                                                       IF (TEST .EQ. 2) GO TO 220
          NO, IT IS A DUMMY ARGUMENT,
          THEN IT MUST BE FOLLOWED BY
          AN INTEGER AND RIGHT BRACKET
                                                       INEG = 2
          CHECK IF -1 FACTOR IS NEEDED
                                                       GO TO 500
                                                       CALL GSCANR($990, IND, IDT, ITC, ICC)
                                                 285
                                                       IF (IND .NE. 1 .OR. IDT .LE. 0) CALL FMLERR($990, ITC, 1, 1)
                                                       I=ILINK1(NP,5,IDT)
                                                       CALL GSCANR($990, IND, IDT, ITC, ICC)
                                                       IF (IND .NE. 4 .OR. IDT .NE. 3) CALL FMLERR($990:ITC:1:1)
                                                       TEST=4
                                                       GO TO 30
    RIGHT BRACKET
     END OF FUNCTION ARGUMENTS
```

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290
                                                   IF (AND(TEST, 113) .NE. 0) CALL FMLERR($990, ITC, 1,1)
                                                   BRT = BRT-1
                                                   IF (BRT .LT. 0) CALL FMLERR($990, ITC, 1,4)
                                                   TEST=4
                                                   CALL STKOUT ($990,17)
                                                   IF (ITYP(NP) .NE. 16) CALL FMLERR($990, ITC, 1, 1)
                                                   I=D(NP)
                                                   J=NP
                                                   NP=LAST(NP)
                                                   CALL RMOVE (J)
                                                   J=ITYP(NP)
                                                   IF (J .LT. 21) CALL FMLERR($990, ITC, 1, 4)
                                                   IF (J .EQ. 24) ITYP(NP)=I+24
                                                   IF ((J .Eq. 24) .AND. ((I+24) .GT. 31))
                                                     CALL FMLERR ($990, D(NP), 1,3)
                                                   IF (J .LT. 24) H2(NP)=I
                                                   CALL STKOUT($990,21)
                                                   GO TO 30
RIGHT PARENTHESIS )
                                                  IF (AND(TEST, 113) .NE. 0) CALL FMLERR($990, ITC, 1, 1)
CHECK IF THIS IS AN END OF SUBSCRIPT
OR THE END OF A SUBEXPRESSION
                                                   PAR = PAR-1
                                                   IF (PAR .LT. 0) CALL FMLERR($990.ITC.1.4)
                                                   CALL STKOUT($990,18)
                                                   IF (ITYP(NP) .NE. 17) CALL FMLERR($990, ITC, 1, 4)
                                                   IF (D(NP) .NE. 0) GO TO 310
     END OF SUBEXPRESSION.
     REMOVE MATCHING '('
                                                   I=NP
                                                   NP=LAST(NP)
                                                   CALL RMOVF1(I)
                                                   TEST=4
                                                   GO TO 30
... END OF A SUBSCRIPT LIST
     CHECK AND PACK SUBSCRIPTS
                                                  TEST=2
                                             310
                                                   N2=D(NP)
                                                   IF (N2 .GT. 4) CALL FMLERR($990,N2,0,5)
                                                   N3=0
                                                   SBC=SBC-1
                                                   IF (SBC .EQ. 0) SB= .FALSE.
                                                   DO 320 KK=N2-1,0,-1
                                                   K=NEXT(NP)
                                                   IF (ITYP(K) .NE. 0) CALL FMLERR($990.D(K).2.13)
                                                   IF (D(K) .LT. 0 .OR. D(K) .GT. 511) CALL FMLERR($990.D(K).0.15)
                                                   FLD(9*KK+9,N3) = D(K)
                                                  CALL RMOVF1(K)
     CHECK IF THIS IS THE END OF
     TRANSLATION (ISW=2)
                                                   IF ((ISW .EG. 2) .AND. (.NOT. SB)) GO TO 330
     :... NO, GO BACK TO VARIABLE PART
TO GET THE VALUE OF THE
          SUBSCRIPTED VARIABLE
                                                   J=NEXT(NP)
                                                   N1=D(J)
                                                   K=NP
```

1 10.4

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NP=LAST(NP)
                                                       CALL RMOVFN(K,J)
                                                       GO TO 140
          :... YES, RETURN FOR ISW=2
                                                 330
                                                       CALL IFREEO(NPO)
                                                       RETURN
:... MINUS
     SET 'NEG= AND LINK IN +
                                                 340
                                                       NEG: .TRUE.
... PLUS
    IS IT UNARY OR BINARY
                                                 350
                                                       IF (AND(TEST, 94) .NE. 0) GO TO 360
     :... UNARY PLUS OR MINUS
                                                       CALL STKOUT ($990,18)
                                                       TEST=16
                                                       GO TO 30
     :... BINARY + -
                                                 360
                                                       J=18
     COMMON PART FOR BINARY OPERATORS
                                                 370
                                                 380
                                                       IF (AND(TEST, 113) .NE. 0) CALL FMLERR($990, ITC, 1, 1)
                                                       TEST=16
                                                       CALL STKOUT($990.J)
                                                       NP=ILINK1(NP,J,I)
                                                       GO TO 30
:... MULTIPLICATION *
                                                 390
                                                       J=19
     GO TO BINARY OPERATOR
                                                       GO TO 370
:... EXPONENTIAL. **
                                                 400
                                                       J=20
     GO TO BINARY OPERATOR
                                                       GO TO 370
:... DIVISION /
                                                 410
                                                       1=-2
                                                       J=19
    GO TO BINARY OPERATOR
     SECOND ENTRY
                                                       GO TO 380
:... LEFT PARENTHESIS (
                                                 420
                                                       IF (AND (TEST: 78) .NE. 0) CALL FMLERR ($990: ITC: 1:1)
                                                       PAR = PAR+1
                                                       TEST=32
                                                       INEG = 3
    GO TO CHECK FOR -1 FACTOR
                                                       GO TO 500
                                                 425
                                                      NP=ILINK1(NP,17.0)
                                                       GO TO 30
:... COMMA
                                                 430 IF (AND (TEST, 113) .NE. 0) CALL FMLERR($990, ITC, 1, 1)
                                                       TEST=32
                                                       CALL STKOUT ($990,18)
                                                       D(NP)=D(NP)+1
                                                       GO TO 30
... EQUAL SIGN
                                                      IF (AND (TEST, 117) .NE. 0) CALL FMLERR ($990, ITC, 1, 1)
                                                       TEST=1
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IF (EOL .OR. SB) CALL FMLERR($990,ITC,1,1)
                                                                EQL= .TRUE.
                                                                CALL STKOUT($990,18)
               GET AND CHECK VARIABLE FOR ASSIGN
               STATEMENT SAVE INFO. IN NN1, NN2, NN3
                                                                IF ((ITYP(NP) .NE. 16) .OR. (D(NP) .NE. 1) .OR. (LAST(NP) .NE. ))
                                                               1 CALL FMLERR($990, ITC, 1, 1)
                                                                KK=NEXT(NP)
                                                                IF (ISW .EQ. 1) GO TO 450
                                                                IF ((KK .EQ. 0) .OR. (ITYP(KK) .LT. 6) .OR. (ITYP(KK) .GT. 7))
                                                               1 CALL FMLERR($990, ITC, 1, 1)
                                                                NN1=D(KK)
                                                                NN2=0
                                                                IF (ITYP(KK) .EQ. 6) GO TO 450
                                                                NN3=NEXT(KK)
                                                                NN2=ITYP(NN3)-7
                                                                NN3=D(NN3)
                                                          450
                                                                CALL IFREEO(KK)
                                                                NEXT(NP)=0
               GO BACK TO TRANSLATE EXPRESSION
                                                                GO TO 30
          :... SEMICOLON ;
          :... APOSTROPHE '
               GO TO END OF TRANSLATION
                                                                IF (AND(TEST:113) .NE. 0) CALL FMLERR($990:ITC:1:1)
END OF LOOP
END OF TRANSLATION
                                                                IF (PAR .NE.O .OR. BRT .NE.O) CALL FMLERR($990, 1() [31,1,4)
                                                                CALL STKOUT ($990, 18)
                                                                IF ((ITYP(NP) .NE. 16) .OR. (LAST(NP) .NE. 0))
                                                                         CALL FMLERR($990, ITC, 1, 1)
    RETURN FOR 'ISW' = 1 AND 3
                                                                IF (ISW .NE. 0) RETURN
                                                                IF (.NOT. EQL) CALL FMLERR($990,ITC,1,1)
    TRANSFER ASSIGNED VARIABLE INFORMATION
                                                                N1=NN1
                                                                N2=NN2
                                                                N3=NN3
     RETURN FOR 'ISW' =0
                                                                RETURN
CHECK IF -1 FACTOR MUST BE INSERTED
     INDICATED BY 'NEG'
     :... NO, GO TO RETURN TO CALLING PLACE
                                                          500
                                                                IF (.NOT. NEG) GO TO 510
     :... YES, LINK IT IN
                                                                NP = ILINK1(NP,19,2)
                                                                CALL ILINK1(NP,0,-1)
                                                                NEG = .FALSE.
          RETURN TO CALLING PLACE
                                                                GO TO (50,285,425), INEG
ERROR RETURN
                                                          990
                                                                CALL IFREEO(NPO)
                                                                RETURN 1
                                                                END
```

p ...4

#### APPENDIX B-PRINTED SUBPROGRAM: EXAMPLE 2

... ERASE STATEMENT

\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\* TITLE MAIN PROGRAM FOR INTERACTIVE FORMAL SYSTEM PARAMETER IDIM = 10" DIMENSION IN(14), INN(14), ITAB(IDIM) EQUIVALENCE (IN(2), INN(1)) DATA INN(14) / ';' /
DATA ITAB /'READ PRINT DUMP ERASE OPTIONCOMMEN' + , 'ROLOUTNCOUNTSAVE RESET '/ \* SEQUENCE CHART INITIALIZE BY CALLING FMLOPT CALL FMLOPT ('INT: ',0) LOOP TO GET NEXT INPUT LINE READ LINE READ 100, END=200, IN FORMAT (13A6, A2) 100 IF IT STARTS WITH 'C ' (COMMENT), GO TO GET NEXT IF (FLD(0,12,IN(1)) .EQ. 1005K) GO TO 110 IF IT STARTS WITH 'P ' (PRINT), GO TO 'P ' ENTRY IF (FLD(0,12,IN(1)) .EQ. 2505K) GO TO 22 J = 0LOOP TO GET STATEMENT TYPE IN J DO 111 I = 1. IDIM IF  $(IN(1) \cdot EQ \cdot ITAB(I)) J = I$ END OF LOOP J=0, IT IS AN ASSIGN STATEMENT IF (J) ,60, REPRINT ERASE, OPTION, ROLOUT, SAVE AND RESET STATEMENTS CONTINUE PRINT 101, IN FORMAT (XA6, 1:1, 13A6) BRANCH BY TYPE GO TO (1: 2: 3:4: 5: 110: 7: 8: 9: 10): J : ... READ STATEMENT CALL FMLIO1 (INN,0) GO TO 110 :... PRINT STATEMENT CALL FMLIO2 (INN.0) GO TO 110 'P ' = PRINT FLD(0.6, IN(1)) = 0505KCALL FML102 (IN, 0) GO TO 110 DUMP STATEMENT

CALL ONDMP

CALL DUMP(K) CALL OFFDMP GO TO 110

CALL FMLERS (INN.0)

IF (INN(1) .NE. ' ') K = 0

```
GO TO 110
      ... OPTION STATEMENT
                                                                                 CALL FMLOPT (INN, 0)
GO TO 110
      ... ROLOUT STATEMENT
                                                                                 CALL FMLOUT (INN,0)
GO TO 110
      ... NCOUNT STATEMENT
                                                                                 CALL COUNT
GO TO 110
      ... SAVE STATEMENT
                                                                                 CALL FMLSAV (INN)
GO TO 110
                                                                          9
      ... RESET STATEMENT
                                                                                 CALL FMLRES (INN)
GO TO 99
                                                                          10
      ASSIGN STATEMENT
                                                                                 PRINT 102, IN
FORMAT (X14A6)
CALL FMLASG (IN,0)
GO TO 110
                                                                          102
END OF FILE READ - STOP
                                                                                 STOP
                                                                          200
                                                                                 END
```

# CHMN PROC

```
***********************************
            TITLE
COMMON DATA STRUCTURE FOR FORMAL SYSTEM
***********************************
            DATA STURERUE
ARRANGED IN 3 LABELED COMMONS
USED AS PROCEDURE. INCLUDED IN OTHER SUBPROGRAMS
                                                                IMPLICIT INTEGER (A-Z)
                                                                PARAMETER ERROR = ERRERR
1. LINKED STORAGE AREA
    THE CORRESPONDING C(1)-D(1) WORDS ARE ALWAYS
     USED IN PAIRS FOR STORING AN ITEM.
     THE DIMENSION OF COD, COIM, MAY BE CHANGED
     DURING INSTALLATION.
     FIELDS IN THE C-D WORDS DEPEND ON THE USAGE.
     THEY ARE DEFINED BY PROCEDURE 'PWORD'. GENERALLY.
     THE LAST 15 BITS IN C IS USED FOR LINKAGE OF
    LINEAR ARRAYS.
                                                                PARAMETER COIH # 2048
                                                                COMMON /FHLCH2/ CICDIMI
                                                                COMMON /FMLCM3/ D(CDIH)
2. COMMON BLOCK FOR INDIVIDUAL POINTERS AND SWITCHES
                                                                COMMON /FMLCH1/
**** FREE (AVAILABLE) STORAGE IN C-D
         C(NXNXNX) = FIRST
         C(ILILIL) - LAST LOCATION
          THE LINEAR ARRAY IS LINKED IN THE
         LAST 15 BITS OF THE C-WORDS.
                                                                    NXNXNX+ ILILIL
... SYMBOL TABLE WITH TREE STRUCTURE IN 4 LEVELS
          STORED IN C-D AREA. FIELDS IN THE C-WORD!
             ITYPB - LAST - NEXT
          NS * FIRST ENTRY IN CINS)-DINS)
                                                                    NS,
          NSB - SUBROUTINE LEVEL POINTER
               SUBPROGRAMS ARE IN ALPHABETIC ORDER
                                                                    NSB .
          I ... ITYPB(NSB) - 0
          1 . . . DINSBI . ALPHANUMERIC NAME OF THE
              SUBROGRAM
          : . . . NSY = LASTINSB) POINTER TO SYMBOL ENTRY
                   SYMBOLS ARE IN ALPHABETIC ORDER
                   NSVI - POINTER TO PRECEEDING SYMBOL
                   ENTRY
                                                                    NSY.NSYI.
              1 ... DINSY) = AUPHANUMERIC NAME OF THE
                   SYMBOL
               **** ITYPBINSY ! TYPE OF SYMBOL,
                        SEE TABLE I.
                    1 ... >31 . INDIRECT REFERENCE
                        LASTINSY POINTS TO AN OTHER
                        SYMBOL
                    *** SECOND AND THIRD BIT * 11
```

SUBSCRIPTED VARIABLE

APPENDIX C-PRINTED SUBPROGRAM: EXAMPLE 3

```
NSU # LAST(NSY). POINTER TO
                              SUBSCRIPT ENTRY
                              SUBSCRIPTS ARE ORDERED BY
                              NUMBER OF SUBSCRIPTS: AND
                              BY ACTUAL SUBSCRIPTS, NSUI
                              - POINTER TO PRECEEDING
                              SUBSCRIPT ENTRY
                                                                     NSU, NSUI,
                         :... b(NSU) = NUMERIC SUBSCRIPT.
                              SEE TABLE 11.
                         1... FIRST 3 BITS OF CINSULE
                              *** = 110. UNASSIGNED.
                                   INDIRECTLY REFERENCED
                                        LASTINSUI POINTS
                                        TO OTHER SYMBOL
                                        ENTRY
                              **** #111, LAST SUBSCRIPT
                                   ENTRY.
                                        LAST (NSU) POINTS
                                        BACK TO 1TS
                                        SYMBOL - NSY
                              1 ... = 010. NORMAL ENTRY
                                        LAST(NSU) POINTS
                                        TO EXPRESSION
                                        VALUE . EPTR
                         1 ... NEXTINSU) . FORWARD LINK TO
                              NEXT SUBSCRIPT.ZERO FOR
                    THE LAST ONE
                         EPTR . EXPRESSION POINTER
                                                                     EPTR.
                         1 ... | TYPB(EPTR) = 16
                         I... LASTIEPTRI = 0
                         : . . . H2(EPTR) = NUMBER OF
                              EXPRESSIONS (FOR LISTS)
                         $ . . . HI (EPTR) =
                             1 . . . = D. EXPRESSION IS IN
                              : CORE
                              LARA NOT ZERO, EXPRESSION
                                        15 04 2844
                                        INDEX . HI (EPTR)
                         ANTA CRANNOS (FIRST TAXES LINK
                              TO THE LINEARLY STORED
                              EXPRESSION WHEN IT IS ON CONFESTIVE
                              EUTRIES ARE ACCORDING
                              . 111 BARAT CT
               * ... NEXTINSY! . FORHARD LINK TO NEXT
                    STHADL, ZERO FOR LAST ONE
          TEAM OF SHILL CHANNON . (EEM) TXAM
               SUBPROSRAM, ZERG FOR LAST ONE
**** TEMPORARY VARIABLES FOR
         NI * NAME OF A VARIABLE
         N2 * NUMBER OF SUBSCRIPTS
         N3 # SUBSCRIPT #ORD
          TO RO + 3 OR I FOR SUBSCR+ OR NOT
```

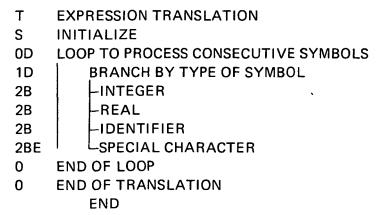
```
N1 . N2 . N3 . 17 .
... OPTION SWITCHES
          XOTOPT - OPTION WORD FROM WAST STATEMENT
          PRODEX . EXPAND PONERS OVER PRODUCT
          INTES* . EVALUATE INTEGER VALUED FUNCTIONS
          MATHSH . EVALUATE MATHEMATICAL FUNCTIONS
          EXPOSM - USE DISTRIBUTIVE LAW
          POWER . EXPAND SUMS RAISED TO POS. INTEGERS
          BASE = 0,1,2,3 FGR BASE(0),(2),(10),(E)
                                                                     XQTOPT . PRODEX . INTGS # . MATHS W . EXPOS W . POWER . BASE .
... MISCELLANECUS
          SIMPSA . USEC BY STOUT ROUTINE FOR RECURSIVE
                   SIMPLIFICATION
          BITSW . USED BY STOUT ROUTINES
          IGUNIT = 1/0 UNIT NUMBER IF 1/0 STATEMENTS
          FTRARG . NUMBER OF FORTRAN TYPE ARGUMENTS
          DEFARE = AUMBER OF ARGUMENTS IN A DEFINED
                   FUNCTION
          DEFFUN . 1 IF CEFINED FUNCTION, O FOR VARIABLE
          NK . START OF ARGUMENT CHAIN IN C-D FOR LIST
               OF VARIABLES
          CBUF # 1/0 BUFFER
          NP . PUSH-DOBN STACK POINTER IN C-D AREA
                                                                     SIMPSW.BITSW. IOUNIT, FTRARG, DEFARG, DEFFUN, NK, CBUF, NP
*********************************
                                                                 LOGICAL INTGS# MATHSW . POWER . SIMPSW . BITSW . PRODEX
                                                                 REFERENCES ON
                                                            END
```

# APPENDIX D-DEFINITION OF INITIAL SEQUENCE CHART

# **Coding Form**

The coding form is divided into three fields: Field 1 consists of one character, the general directive for input; field 2 contains special directives for flowchart elements and a label for program statements; field 3 contains the text.

An initial program is illustrated below:



# Input Form

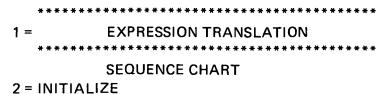
The actual input does not contain the lines; the text is left adjusted in field 3:

Т **EXPRESSION TRANSLATION** S INITIALIZE LOOP TO PROCESS CONSECUTIVE SYMBOLS 0D **BRANCH BY TYPE OF SYMBOL** 1D 2B INTEGER 2B **REAL** 2B **IDENTIFIER** 2BE SPECIAL CHARACTERS **END OF LOOP** 0 **END OF TRANSLATION** 

# **Output Form**

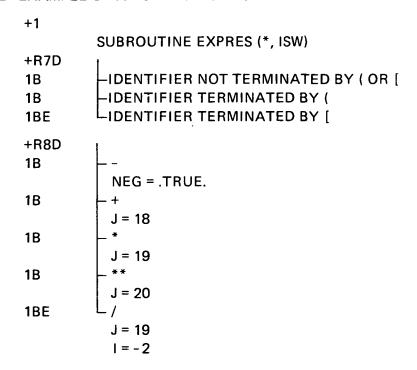
The initial program can be listed with line numbers as follows:

**END** 



```
3 = LOOP TO PROCESS CONSECUTIVE SYMBOLS
4 =: BRANCH BY TYPE OF SYMBOL
5 =: :...INTEGER
6 =: :...REAL
7 =: :...IDENTIFIER
8 =: :...SPECIAL CHARACTERS
9 = END OF LOOP
10 = END OF TRANSLATION
11 = END
```

## APPENDIX E-EXAMPLE OF AN UPDATING PROCEDURE



Note that the '+' is an insertion directive. The number following + indicates the line where the insertion is to be done. 'R' indicates that the levels of lines following to be inserted are defined relative to the line where the insertion occurs.

#### **DISCUSSION**

MEMBER OF THE AUDIENCE: I notice that you have many comments noted through there. It seems to be about a two-to-one comment per statement. Is that about correct?

**MESZTENYI:** It depends on the program. It depends on the language, too. The comments should be semantic, not repeated as an equation.

MEMBER OF THE AUDIENCE: Do you think that some of the discussions about what we can get out of the compiler would fall into this?

MESZTENYI: I would like to have the compiler in the subroutine. I would like to do

that, but I would start here from the development point first, because this is where one defines the program first.

MEMBER OF THE AUDIENCE: It seems that the compiler could give you certain information, and you could add some personal comments and have better descriptive material. Is that true?

MESZTENYI: It depends on what standpoint you look at. As I look at it, I want an overall view from the beginning. Before I finish the program, I might want to give the specification a bigger flowchart type of definition that could be used right away.

MEMBER OF THE AUDIENCE: You are trying to get the flavor of the program that you are working on for a certain purpose. The compiler will only come out with standard words for any program. The compiler does not know what your program is, but you do. With personal comments added to the program, what you have would provide additional information.

MESZTENYI: I find it is hard for programmers to add something after they have written the program. When they write, they do not mind writing down their comments.

MEMBER OF THE AUDIENCE: I am working from the viewpoint that we now have difficulty at times getting any comments in, and if we provided a lead into the comments and they went down the list and it did not make too much sense to them from a general viewpoint, that they could add these rather well.

MESZTENYI: I agree that they could, and this is actually what is now done. I added this myself.

The other part I would like to focus on a little bit is the programming part. If you start from the sketch with those lines coming down and write, you make the programmer apply a little discipline to the subject of program placement. For example, I try to avoid any GO TO unless it is some kind of loop structure. I try to avoid going back. I find a loop for each logic curve that I process, but it is not a DO statement, and I jump directly back to the beginning. It probably would have been much nicer documenting it to go to the end of this loop and comment it, which goes back and gets the next one. In this way it forces the programmer to do a documented description because it is very hard to document a graph that points out the actual information. The text or the description of the program is sequential, but semantically it is a graph. A tree, which is sort of in-between, is much easier to represent. You have cross-references, but the form is still a tree, and this is what I tried to simulate.

MEMBER OF THE AUDIENCE: I think the speaker is trying to get the programmer to write down what is being accomplished and when. Once in the right-hand side, the language does not really matter. He is trying to read narrative text so that you get some concept of when things happen and what really is happening because the specification of the problem is written in a narrative form. He does that rather than deduce what was done from how something is being done. I do not think a programmer is going to do that very well because he is so involved in the mechanics that he cannot get out of them.

MEMBER OF THE AUDIENCE: It seems to me that here is a case where we can go from the rationale of a subroutine and in an automated way feed in the programming language statements. Is this what you had in mind? I could see how you actually tried to develop your subroutine. I can see how you can start with the rationale of the subroutine

first and then by using the type of coding that you did, you could automatically call for the appropriate programming language statements.

**MESZTENYI**: Not automatically. I certainly think of more than just the semantic type of description that I want to accomplish. What I want to accomplish eventually is the statements.